A Few Good Examples why to Refute Classical Logic in Quantum Mechanics

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- Classical logic is a type of formal logic which is most intensively studied and most widely used on all levels of education.
- Classical logic is defined by a set of properties.
- Classical logic generally includes propositional logic and first-order logic.

Definition of CL

We define classical logic with the following set of properties:

- Binarity: Set of truth values - V(P) ∈ {0,1}
- Commutativity: $P \land Q \equiv Q \land P$, $P \lor Q \equiv Q \lor P$
- Distributivity:
 - $P \land (Q \lor R) \equiv (P \land Q) \lor (P \land R)$ $P \lor (Q \land R) \equiv (P \lor Q) \land (P \lor R)$
- Principle of excluded middle: $\mathcal{V}(P \lor \neg P) = 1$
- Principle of non-contradiction: $\mathcal{V}(\neg (P \land \neg P)) = 1$, alternatively $\mathcal{V}(P \land \neg P) = 0$

- Branch of physics
- Atomar and subatomar scale
- Commutative and non-commutative operators
- Uncertainty principle

- In 1936 Birkoff and von Neumann wrote the article "The Logic of Quantum Mechanics".
- Birkoff and von Neumann wanted to find the logical structure in quantum mechanics which did not conform to classical logic.

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Should we refute classical logic for QM?

Definitions

- Quantum mechanical proposition is a projection operator (generally observable) *P*
 - experiment outcome prediction

•
$$P_X = x = 1 \cdot x$$
 for $x \in L_P$
 $P_Y = 0 = 0 \cdot y$ for $y \in L_P^{\perp}$

- $P \land Q$ operator projecting on $L_{P \land Q} = L_P \cap L_Q$
- $P \lor Q$ operator projecting on $\overline{L_P + L_Q}$
- $\neg \mathsf{P} = 1 \mathsf{P}$
- $P \wedge Q = \lim_{n \to \infty} (PQ)^n$
- $\mathbf{P} \lor \mathbf{Q} = \neg [\neg \mathbf{P} \land \neg \mathbf{Q}]$

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Examples

• Example 1:
$$S_z = \left(\frac{\hbar}{2}\right)P_1 - \left(\frac{\hbar}{2}\right)P_2$$

 $P_1 = |\uparrow\rangle \langle\uparrow| \quad P_2 = |\downarrow\rangle \langle\downarrow|$
Matrix notation: $|\uparrow\rangle = \begin{pmatrix}1\\0\end{pmatrix} \quad |\downarrow\rangle = \begin{pmatrix}0\\1\\\end{pmatrix}$
 $P_1 = \begin{pmatrix}1 & 0\\0 & 0\end{pmatrix} \quad P_2 = \begin{pmatrix}0 & 0\\0 & 1\end{pmatrix}$
For S_z :
 $S_z = \frac{\hbar}{2} \begin{pmatrix}1 & 0\\0 & -1\end{pmatrix}$

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Using previous definitions:

•
$$\mathbf{P}_1 \wedge \mathbf{P}_2 = \lim_{n \to +\infty} (\mathbf{P}_1 \mathbf{P}_2)^n$$

$$= \lim_{n \to +\infty} \left[\begin{pmatrix} 1 & 0 \\ 0 & 0 \end{pmatrix} \cdot \begin{pmatrix} 0 & 0 \\ 0 & 1 \end{pmatrix} \right]^n$$

$$= \lim_{n \to +\infty} \begin{pmatrix} 0 & 0 \\ 0 & 0 \end{pmatrix}^n = \mathbf{0}$$

•
$$P_1 \lor P_2 = 1 - [(1 - P_1) \land (1 - P_2)]$$

= $1 - [P_2 \land P_1]$
= 1

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Calculations were simplified due to $[P_1, P_2] = 0$

Example of non-commuting observables:

•
$$S_x = \begin{pmatrix} \frac{h}{2} \end{pmatrix} Q_1 - \begin{pmatrix} \frac{h}{2} \end{pmatrix} Q_2$$
, it is easily seen that:
 $Q_1 = \frac{1}{2} \begin{pmatrix} 1 & 1 \\ 1 & 1 \end{pmatrix} \quad Q_2 = \frac{1}{2} \begin{pmatrix} 1 & -1 \\ -1 & 1 \end{pmatrix}$
are projectors projecting on subspaces $\begin{pmatrix} 1 \\ 1 \end{pmatrix}$, $\begin{pmatrix} 1 \\ -1 \end{pmatrix}$
• Operators Q_1 and Q_2 do not commute with P_1 and P_2
• $\mathbf{P}_1 \wedge \mathbf{Q}_1 = \lim_{n \to +\infty} (P_1 Q_1)^n$
 $= \lim_{n \to +\infty} (\frac{1}{2})^n \begin{pmatrix} 1 & 1 \\ 0 & 0 \end{pmatrix} = \mathbf{0}$
• $\mathbf{P}_1 \vee \mathbf{Q}_1 = 1 - ((1 - P_1) \wedge (1 - Q_1))$
 $= 1 - (P_2 \wedge Q_2)$
 $= 1 - 0 - \mathbf{1}$

Non-commutativity \neq zero conjunction

Example: particle with spin 1

•
$$B = \frac{1}{2\hbar^2} [S^2 - S_z^2] + \frac{1}{2\hbar^2} S_z = \begin{pmatrix} 1 & 0 & 0 \\ 0 & 1 & 0 \\ 0 & 0 & 0 \end{pmatrix}$$

• $D = \frac{1}{\hbar^2} S_z^2 = \begin{pmatrix} \frac{1}{2} & 0 & \frac{1}{2} \\ 0 & 1 & 0 \\ \frac{1}{2} & 0 & \frac{1}{2} \end{pmatrix}$, $[B, D] \neq 0$ but $\begin{pmatrix} 0 \\ 1 \\ 0 \end{pmatrix}$ is a common eigenvector.

,

•
$$D \wedge B = \lim_{n \to +\infty} \begin{pmatrix} \frac{1}{2} & 0 & \frac{1}{2} \\ 0 & 1 & 0 \\ \frac{1}{2} & 0 & \frac{1}{2} \end{pmatrix}^n = \begin{pmatrix} 0 & 0 & 0 \\ 0 & 1 & 0 \\ 0 & 0 & 0 \end{pmatrix}$$

• Easily calculated as before:

$$P_1 \lor P_2 = 1 \ Q_1 \land P_1 = 0 \ Q_1 \land P_2 = 0$$

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• Easily calculated as before: $P_1 \lor P_2 = 1$ $Q_1 \land P_1 = 0$ $Q_1 \land P_2 = 0$

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$$Q_1 \wedge (P_1 \vee P_2) \equiv (Q_1 \wedge P_1) \vee (Q_1 \wedge P_2)$$

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$$Q_1 \wedge 1 \equiv 0 \lor 0$$

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Comment: Modular law holds:

• $C \rightarrow A$

•
$$A \land (B \lor C) \equiv (A \land B) \lor (A \land C) \equiv (A \land B) \lor C$$

• let
$$P_{\epsilon} = \begin{pmatrix} \cos^2 \epsilon & \cos \epsilon \sin \epsilon \\ \cos \epsilon \sin \epsilon & \sin^2 \epsilon \end{pmatrix}$$

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• $\lim_{\epsilon \to 0} P_{\epsilon} = P_1$

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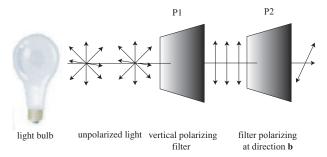
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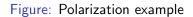
•
$$(\lim_{\epsilon \to 0} P_{\epsilon}) \wedge P_1 = P_1 \wedge P_1 = P_1$$

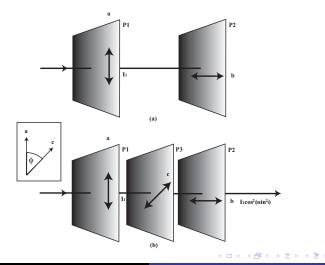
Figure: Polarization example



 Polarizing filter = Mechanical analogue to the projection operator

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• H - horizontal filter, V - vertical filter, D - diagonal filter

•
$$H^n = H$$
, $V^n = V$, $HV = 0$

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• $VHD \neq VDH$

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$$H^n = H, V^n = V, HV = 0$$

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•
$$H \wedge V = \lim_{n \to \infty} (HV)^n = 0$$

•
$$V \wedge D = \lim_{n \to \infty} (VD)^n = 0$$

•
$$D \lor V = \neg [\neg D \land \neg V]$$

= $\neg [D_{anti} \land H] = \neg 0 = 1$

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Failure of distributivity

•
$$H \land (D \lor V) \equiv H \land 1 \equiv H$$

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Failure of distributivity

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$$H \land (D \lor V) \equiv H \land 1 \equiv H$$

•
$$(H \land D) \lor (H \land V) \equiv 0 \lor 0 \equiv 0$$

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• D_{ϵ} - polarizing filter inclined at ϵ to the horizontal $\lim_{\epsilon \to 0} D_{\epsilon} = H$

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- $D_{\epsilon} \wedge H = 0$
- $\lim_{\epsilon \to 0} (D_{\epsilon} \wedge H) = 0$
- $(\lim_{\epsilon \to 0} D_{\epsilon}) \wedge H = H$

Two slit experiment - failure of distributivity

- A₁ electron passes through slit 1 A₂ - electron passes through slit 2
- P(A₁, R) probablity that electron hits region R, after A₁
 P(A₂, R) probablity that electron hits region R, after A₂
- Classically, one would expect $P(R) = \frac{1}{2}P(A_1, R) + \frac{1}{2}P(A_1, R)$

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- Classically, one would expect $P(R) = \frac{1}{2}P(A_1, R) + \frac{1}{2}P(A_1, R)$
- Quantum formula $P(R) = \frac{1}{2}P(A_1, R) + \frac{1}{2}P(A_1, R) + \sqrt{P(A_1, R)P(A_2, R)} \cos \delta$

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•
$$P(A_1 \lor A_2, R) = \frac{P((A_1 \lor A_2) \cdot R)}{P(A_1 \lor A_2)}$$

= $\frac{P((A_1 \cdot R) \lor (A_2 \lor R))}{P(A_1 \lor A_2)}$
= $\frac{P(A_1 \cdot R)}{P(A_1 \lor A_2)} + \frac{P(A_2 \cdot R)}{P(A_1 \lor A_2)}$

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•
$$P(A_1) = P(A_2)$$

 $P(A_1 \lor A_2) = 2P(A_1) = 2P(A_2)$

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 $= \frac{P((A_1 \lor R) \lor (A_2 \lor R))}{P(A_1 \lor A_2)}$
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• $\frac{P(A_1 \cdot R)}{P(A_1 \lor A_2)} = \frac{P(A_1 \cdot R)}{2P(A_1)} = \frac{1}{2}P(A_1, R)$
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$$P(A_1 \lor A_2, R) = \frac{P((A_1 \lor A_2) \cdot R)}{P(A_1 \lor A_2)}$$

 $= \frac{P((A_1 \lor R) \lor (A_2 \lor R))}{P(A_1 \lor A_2)}$
 $= \frac{P(A_1 \cdot R)}{P(A_1 \lor A_2)} + \frac{P(A_2 \cdot R)}{P(A_1 \lor A_2)}$
• $P(A_1) = P(A_2)$
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• $P(R) = \frac{1}{2}P(A_1, R) + \frac{1}{2}P(A_1, R)$

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The question was:

Should we refute classical logic for QM?

And the anwser is:

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Should we refute classical logic for QM?

And the anwser is:

We should.

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